An Attribute-Based Encryption Scheme Based on Unrecognizable Trapdoors

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Abstract: Attribute-Based Encryption (ABE) has been widely used for ciphertext retrieval in the cloud environment. However, bi-flexible attribute control and privacy keywords are difficult problems that have yet to be solved. In this paper, we introduce the denial of access policy and the mutual matching algorithm of a dataset used to realize bidirectional control of attributes in the cloud server. To solve the problem of keyword privacy, we construct a security trapdoor by adding random numbers that effectively resist keyword guessing attacks from cloud servers and external attackers. System security is reduced to the Deterministic Bilinear Diffie-Hellman (DBDH) hypothesis problem. We validate our scheme through theoretical security analysis and experimental verification. Experiments are conducted on a real dataset, and results show that the scheme has higher security and retrieval efficiency than previous methods.

Key words: Attribute-Based Encryption (ABE); unrecognizable trapdoor; two-way access strategy; ciphertext search

1 Introduction

Cloud storage is a mainstream online storage method that removes the hardware and management overhead of users’ local storage and takes physical data out of users’ control. In this environment, data security is greatly threatened. Data encryption is generally adopted to solve the data security problem in cloud storage, but the availability of encrypted data is limited, especially because the widely used keyword retrieval technology of plaintext information cannot be directly applied to encrypted data. Searchable encryption is a good solution for ciphertext retrieval and has three main research directions: security, accuracy, and efficiency.

1.1 Related work

To enhance the security of data on a server, Song et al. proposed a searchable encryption scheme with a one-to-one mechanism. However, the multi-user environment is challenging for searchable encryption. Boneh et al. proposed a searchable encryption scheme based on public key cryptography and proved that the public key searchable encryption system satisfies semantic security, but the scheme cannot resist keyword guessing attacks. A keyword guessing attack is an attack mode in searchable encryption that mainly involves statistical analysis of the uploaded trapdoors. The keywords generate fixed trapdoors and can thus be detected by the server or the attacker during mass uploads. Fang et al. proposed a public key encryption scheme that resists keyword guessing attacks without random prediction, but server internal attacks are the weakness of this scheme. Shao and Yang also proposed a scheme that can resist server keyword guessing attack. The best defense against keyword guessing attacks is to make the trapdoors unrecognizable, which means that the same keyword produces a different trapdoor each time. Privacy protection of keywords in searchable encryption mainly ensures the privacy protection of trapdoors. How to construct secure trapdoors is a difficult problem in searchable encryption.

To make ciphertext retrieval more flexible and
efficient, Attribute-Based Encryption (ABE)\textsuperscript{[3]} is proposed. Among them, the Ciphertext-Policy Attribute-Based Encryption (CP-ABE)\textsuperscript{[6]} can embed the access control strategy in ciphertext and control the access of user attributes flexibly, but introduce the problem of attribute cancellation. Data in the cloud environment are dynamic and massive. Hence, the attribute-based encryption scheme needs to modify attributes and policies efficiently. Pirretti et al.\textsuperscript{[7]} proposed a secure attribute-based system scheme that could revoke attributes. By setting an expiration date for attributes, the authorizing agency periodically updates the attribute version and revokes user attributes by revoking the latest version of an attribute. However, attribute-based encryption schemes have security, accuracy, and efficiency problems. To solve the security problem, Hur and Noh\textsuperscript{[8]} proposed attribute-based encryption scheme with attributes and user revocation capabilities, which enhances the forward and backward security of user access control and has attribute revocation ability. Li et al.\textsuperscript{[9]} proposed a searchable ciphertext policy attribute-based encryption scheme, which has the ability to revoke attributes in the Cloud Server (CS). To protect the privacy of users, Ma et al.\textsuperscript{[10]} proposed a privacy-preserving multi-authority CP-ABE with a revocation scheme based on privacy protection, which not only revokes attributes but also protects users’ privacy effectively.

To improve the accuracy of attribute revocation in searchable encryption schemes, Yang et al.\textsuperscript{[11]} proposed attribute-based fine-grained access control with an efficient revocation scheme for cloud storage systems. This scheme does not require any cooperative access control by the server, and the Data Owner (DO) does not need to be online in real time. However, the program proves its security under random prediction only. Zu et al.\textsuperscript{[12]} proposed a new CP-ABE with an efficient revocation scheme. The access structure of the scheme adopts the Linear Secret Sharing Scheme (LSSS) mode and has strong performance capability. For more granular access, Sun et al.\textsuperscript{[13]} used CP-ABE and proxy re-encryption to implement file-level access authorization and support data User (U) attribute revocation. Cui et al.\textsuperscript{[14]} proposed ABE with an expressive and authorized keyword search scheme that is more accurate and achieves fine-grained access control of encrypted data in the cloud.

In terms of revocation efficiency, Xue et al.\textsuperscript{[15]} proposed a ciphertext comparable attribute-based encryption scheme based on 0-1 encoding and an efficient construction method based on the generation and management of subattributes with 0 and 1 encoding concepts to reduce communication and computational overhead. Outsourcing computing can greatly reduce user online computing cost. Hence, Chen et al.\textsuperscript{[16]} proposed an online-offline ciphertext policy attribute-based searchable encryption scheme that uses offline preprocessing outsourcing decryption and reduces user online computing cost while improving efficiency. To reduce the computational cost of outsourcing decryption, Zhao et al.\textsuperscript{[17]} proposed a constant cipher-sized attribute-based encryption scheme. The scheme’s ciphertext size is constant, which not only improves the outsourcing computational efficiency but also makes the system efficient.

Research on ciphertext retrieval has become increasingly diversified. Qian et al.\textsuperscript{[18]} proposed a ciphertext policy attribute-based searchable encryption scheme for multi-authorization centers in the cloud environment. The scheme uses re-encryption technology to update the ciphertext during the process of attribute revocation, and a multi-authorization center effectively improves the overall efficiency of the program. Broadcast encryption is a typical one-to-many mode. Canard et al.\textsuperscript{[19]} combined broadcast encryption with attribute-based encryption to form a new secret sharing method for one-to-many searchable encryption modes. Xue et al.\textsuperscript{[20]} combined deterministic deletion with attribute-based encryption to propose an attribute-based ciphertext retrieval scheme that supports revocation.

1.2 Our contribution

To solve the problem of trapdoor security and attribute fine-grained access in existing schemes, we propose an Attribute-Based Encryption scheme based on Unrecognizable trapdoors (U-ABE). Our scheme implements one-time trapdoor construction and introduces a denial of access policy that makes attribute control more flexible. The main contributions of this work are as follows.

1. To solve the trapdoor safety problem, we use the bilinearity of bilinear mapping to construct a one-time trapdoor and prove that the trapdoor is unidentifiable.

2. To improve the flexible control of access policy, we introduce a denial of access policy that is controlled in both directions.

3. Through theoretical security analysis, our scheme satisfies the indistinguishable Keywords Guessing Attack (called IND-KGA) and the statistically
indistinguishability under Chosen Ciphertext Attack (called IND-CCA) secure in the Decisional Bilinear Diffie-Hellman (DBDH).

1.3 Organization of this article

The main parts of this paper are organized as follows: Section 2 presents relevant basic knowledge. Section 3 briefly introduces the U-ABE scheme. Section 4 introduces the U-ABE scheme in detail. Section 5 presents an analysis of the accuracy and security of the scheme. Section 6 provides the theoretical and experimental analyses of the scheme. Section 7 summarizes this article.

2 Preliminary

2.1 Bilinear pairing

It is assumed that groups $G$ and $G_T$ are cyclic groups whose order is prime $p$. $g$ is the generator of group $G$, and there exists a bilinear map $\hat{e} : G \times G \rightarrow G_T$, that satisfies the following properties:

1. **Bilinearity:** For any $x, y \in G$ and $a, b \in \mathbb{Z}_p$, we have $\hat{e}(x^a, y^b) = \hat{e}(x^b, y^a) = \hat{e}(x, y)^{ab}$, where $\mathbb{Z}_p$ is the set of nonnegative integers less than $p$.

2. **Non-degeneracy:** $\hat{e}(g, g) \neq 1$.

3. **Computability:** There is an efficient algorithm to compute $\hat{e}(x, y)$ for any $x, y \in G$.

2.2 DBDH

Let groups $G_1$ and $G_2$ and map $\hat{e} : G_1 \times G_1 \rightarrow G_2$. $g$ is the generator of group $G_1$, randomly generate $(a, b, c, z) \leftarrow \mathbb{Z}_p$, and generate two quintuples $T_0 = (g, A = g^a, B = g^b, C = g^c, Z = \hat{e}(g, g)^z)$ and $T_1 = (g, A = g^a, B = g^b, C = g^c, Z = \hat{e}(g, g)^{abc})$. Write the two five-tuples as

$$P_{DBDH} = (g, g^a, g^b, g^c, \hat{e}(g, g)^{abc}),$$

$$R_{DBDH} = (g, g^a, g^b, g^c, \hat{e}(g, g)^{c}).$$

**DBDH assumption:** There is an adversary without polynomial time, who can at least distinguish the quintuples $P_{DBDH}$ and $R_{DBDH}$ with a non-negligible advantage $\varepsilon$.

2.3 List of symbols in the text

The meanings of the symbols in the system are as follows:

- $B$: Encryption parameter
- $B_i$: Decryption parameter
- $t_1$: Allow access policy
- $t_2$: Denial of access policy
- $U_1$: Encrypted allowed access policy
- $U_2$: Encrypted denial of access policy
- $CT$: Ciphertext, including $C_m$ and $C_k$
- $V$: Version information of table
- $\varphi$: Keyword index
- $w$: Data keyword collection
- $w_i$: Retrieve keyword collection
- $S_k_a$: Attribute private key set
- $T_w$: Keyword index
- $p$: Matching information

2.4 Security model

The security model includes the statistically IND-KGA secure and IND-CCA secure.

**Definition 1** The statistically IND-KGA secure allows Adversary $A$ to execute the keyword guessing attack to distinguish the trapdoors corresponding to $w_0$ and $w_1$. Adversary $A$ and Challenger $C$ perform the game as follows.

**Setup:** Given the security parameter $\lambda$, Challenger $C$ runs the initialization algorithm to generate the public parameter $par$.

**Phase 1:** Adversary $A$ runs the Trap$(w, par)$ algorithm multiple times.

**Challenge:** Adversary $A$ randomly selects two keywords $w_0$ and $w_1$ from the keyword space, then sends them to Challenger $C$. Challenger $C$ flips a random coin $\mu \in \{0, 1\}$, runs the algorithm Trap($w$, $par$), and finally sends the trapdoor $T_{w_0}$ to Adversary $A$.

**Phase 2:** Same as Phase 1.

**Guess:** Adversary $A$ outputs a keyword $\mu'$, and if $\mu = \mu'$, Adversary $A$ wins the security game.

**Definition 2** The security model includes the statistically IND-KGA secure and IND-CCA secure.

**Setup:** Given security parameter $\lambda$, Challenger $C$ executes the setup algorithm $Init(\lambda^2)$ to generate the public parameter $par$. Given groups $G_1$ and $G_2$ and map $\hat{e} : G_1 \times G_1 \rightarrow G_2$. Challenger $C$ randomly generates $(a, b, c, z) \leftarrow \mathbb{Z}_p$ to generate a quintuple $T_0 = (g, A = g^a, B = g^b, C = g^c, Z = \hat{e}(g, g)^z)$.\n
**Phase 1:** Adversary $A$ runs the encryption algorithm multiple times.

**Challenge:** Adversary $A$ randomly selects two keywords $m_0$ and $m_1$ from the plaintext space, then sends them to Challenger $C$. Challenger $C$ flips a random coin $\mu \in \{0, 1\}$, runs the encryption algorithm, and finally sends the ciphertext $CT$ to Adversary $A$.

**Phase 2:** Same as Phase 1.

**Guess:** Adversary $A$ outputs a plaintext $\mu$, if $\mu = \mu'$,
Adversary A wins the security game.

3 System Model

The system model of the scheme is shown in Fig. 1. The scheme includes four entities: DO, CS, U, and AA.

1. AA: Assuming that the attribute authority is trusted, its main tasks are to generate a random table of attributes, encrypt the policy uploaded by the DO, and calculate the encryption parameters. The attributes uploaded by the U are calculated according to the table, and the attribute private key and decryption parameters are obtained.

2. DO: The main task of the DO is to encrypt the data by using traditional symmetric encryption. The addition of random numbers generates an unrecognizable index, and the data authority interacts with the AA to obtain an access strategy.

3. U: The main tasks of U are to generate an unrecognizable random trapdoor, perform data interaction with the CS, obtain the version number returned by the server, and upload the version number and its own attributes to the AA. The AA then calculates and returns the private attribute key and decrypts the parameter. Then, the private attribute key is uploaded to the CS and the ciphertext is returned after the server verifies.

4. CS: The main tasks of the CS are to receive the trapdoor uploaded by the U, deliver the ciphertext version number to the U, receive the attribute private key uploaded by the U, and deliver the ciphertext to the U after the matching operation.

The following is an introduction to the algorithms used in this article. There are a total of seven algorithms.

1. Setup($\lambda$) → (Par, T): The trusted AA runs the algorithm, inputs the security parameter $\lambda$, and outputs the public parameter par and the random attribute table. The random attribute table is privately owned and regularly updated by the AA.

2. EncT($t_1$, $t_2$, ss) → ($U_1$, $U_2$, B): The algorithm inputs $t_1$, $t_2$, and ss, and outputs $U_1$, $U_2$, and B. $t_1$ is the set of allowed access policies uploaded by the DO to the AA. $t_2$ is the set of denial of access policies uploaded by the DO to the AA. ss is the coordinates where attributes in the property authority query $t_1$ and $t_2$ set are positioned in the attribute table. $U_1$ and $U_2$ are encrypted from $t_1$ and $t_2$, respectively. After ss encryption, the encryption parameter B is obtained, and the version information V of the table is embedded into encryption parameter B. The AA returns the ciphertext to the DO.

3. Enc($m$, k, w, B, par) → (CT, $\varphi$): The DO calculates the algorithm. $m$, k, w, B, and par are inputted. CT and $\varphi$ are outputted. $m$ is plaintext data, k is a key symmetrically encrypted for plaintext, w is the set of keywords contained in the data, and B is the encryption parameter. The ciphertext CT includes two ciphertexts: one is $C_m$, which is obtained by symmetrically encrypting the plaintext $m$, and the other is $C_k$, which is obtained by encrypting the key k. $\varphi$ is an index obtained by encrypting the keyword set w.

4. Trap(w, par) → $T_w$: The user calculates the algorithm. The w and par are inputted, and the $T_w$ is outputted. w is the keyword when the user queries. After the calculation, the keyword trapdoor $T_w$ is obtained and uploaded to the CS for retrieval.

5. KeyGen(att, V) → SKa: The AA calculates the algorithm. att and V are inputted, and SKa is outputted. att is the attribute uploaded by the user, V is the version information, and SKa is the attribute parameter.

6. Search($U_1$, $U_2$, $\varphi$, $T_w$, SKa) → 1 or 0: The server runs the algorithm for matching retrieval. The algorithm is divided into two parts in the system. The first stage
is to perform a keyword search to obtain version information \( V \). The second stage is to match the attribute with the access policy. The attribute must not intersect with the attribute set in the denied access policy, and it also contains the set of attributes in the allowed access policy. When both stages meet the requirements, the server sends a ciphertext \( CT \) to the user.

(7) \( \text{Dec}(CT, p, B_i) \rightarrow m \): The user calculates the algorithm, inputs ciphertext \( CT \), matches information \( p \), and decrypts parameter \( B_i \). The plaintext \( m \) is obtained by calculation.

## 4 System Specification

\( \text{Setup}(1^k) \rightarrow (Par, T) \): Given the security parameter \( \lambda \), the trusted AA running the algorithm outputs the bilinear mapping \( \text{par} \), random numbers \( (a, b) \in \mathbb{Z}_q \) are generated, and a hash function \( H : (0, 1)^* \rightarrow G_1 \) is defined. Finally, the public parameter \( \text{par} \) is outputted,

\[ \text{par} = (a, b, g, g^a, g^b, G_1, G_2, e, g, H). \]

The AA generates an attribute table, which is private to the AA, and various attributes are placed in the table, as shown in Table 1. Its coordinates represent this attribute in the calculation. The coordinates are composed of random numbers, and the coordinate data are periodically replaced. When the coordinate data in Table 1 are replaced, the AA calculates the new version information \( V \).

\( \text{EncT}(t_1, t_2, ss) \rightarrow (U_1, U_2, B) \): The AA receives the allowed access policy \( t_1 \) and the denial of access policy \( t_2 \) uploaded by the DO, and then the AA searches in Table 1 to obtain the coordinate data \( ss \),

\[ ss = [(x_1, y_1), (x_2, y_2), \ldots, (x_n, y_n)], \]

\[ (x_1, y_1), (x_2, y_2), \ldots, (x_n, y_n),\ 1 \leq i, u \leq n. \]

The AA calculates \( U_1, U_2 \), and encryption parameter \( B \) using the coordinate data in \( ss \),

\[ U_1 = [(g^{x_1}, g^{y_1}), (g^{x_2}, g^{y_2}), \ldots, (g^{x_n}, g^{y_n})], \]

\[ U_2 = [(g^{x_1}, g^{y_1}), (g^{x_2}, g^{y_2}), \ldots, (g^{x_n}, g^{y_n})]. \]

\[ B = \prod_{i=1}^{n} H(x_i^{y_i}). \]

\[ \text{Enc}(m, k, w, B, \text{par}) \rightarrow (CT, \varphi) \]: The ciphertext \( CT \) contains ciphertext \( C_l \) and \( C_m \), \( CT = (C_k, C_m) \), and the DO uses symmetric encryption to encrypt plaintext \( m \) to obtain ciphertext \( C_m \). The encryption key is \( k \), key \( k \) is encrypted, a random number \( t \leftarrow \mathbb{Z}_p \) is generated, and then calculate

\[ C_k = [\hat{e}(g^a, g^b)' \times k, g^t, B^t]. \]

The DO needs to establish an index in the data. Where the set of keywords is \( w \). Then, the random number set \( S \leftarrow \mathbb{Z}_p \) is generated, and then calculate

\[ \varphi = [(g^{x_1}, w_1^{x_1}), (g^{x_2}, w_2^{x_2}), \ldots, (g^{x_n}, w_n^{x_n})]. \]

\[ \text{Trap}(w, \text{par}) \rightarrow T_w \]: The algorithm inputs the keyword set \( w \) that needs to be inquired, generates the random number set \( L \leftarrow \mathbb{Z}_p \), and calculates:

\[ T_w = [(g^{l_1}, w_1^{l_1}), (g^{l_2}, w_2^{l_2}), \ldots, (g^{l_n}, w_n^{l_n})]. \]

\[ \text{KeyGen}(\text{att}, V) \rightarrow \text{SK}_a \]: After data interaction between the user and server, the user receives the version information \( V \) sent by the CS, and the user sends the \( \text{att} \) together with the version information \( V \) to the AA. The AA searches the corresponding attribute data in Table 1 of the corresponding version \( V \) of the user attribute \( \text{att} \), and then calculates the attribute data \( \text{SK}_a \) from the coordinate data,

\[ \text{SK}_a = [(g^{a_1}, g^{b_1}), (g^{a_2}, g^{b_2}), \ldots, (g^{a_n}, g^{b_n})]. \]

\[ \text{Search}(U_1, U_2, T_w, \text{SK}_a) \rightarrow 1 \text{ or } 0, \text{ the retrieval is divided into two phases}. \]

**Phase 1**: The user uploads trapdoor \( T_w \) to the CS, and the server matches the trapdoor with the index. The calculation process is as follows:

\[ \hat{e}(g^t, w_1^{x_1}) = \hat{e}(g^{x_1}, w_1^{l_1}). \]

If the equation is true, the proof is the same keyword. If the equation is false, the next keyword is evaluated until the last keyword in the index. The server sends the version information \( V \) of the data to the user.

**Phase 2**: The user uploads the attribute private key \( \text{SK}_a \) to the server, and the CS matches the attribute with the access policy. Given that two access policies exist in this solution, the attributes contained in the denial of access policy must not be included by the user attribute private key. The user attribute private key \( \text{SK}_a \) needs to be matched with the denial of access policy \( U_2 \) first, and the following is calculated:

<table>
<thead>
<tr>
<th>( x )</th>
<th>( y )</th>
</tr>
</thead>
<tbody>
<tr>
<td>( x_1 )</td>
<td>( y_1 )</td>
</tr>
<tr>
<td>( x_2 )</td>
<td>( y_2 )</td>
</tr>
<tr>
<td>( x_3 )</td>
<td>( y_3 )</td>
</tr>
<tr>
<td>( x_4 )</td>
<td>( y_4 )</td>
</tr>
<tr>
<td>( x_5 )</td>
<td>( y_5 )</td>
</tr>
<tr>
<td>( \vdots )</td>
<td>( \vdots )</td>
</tr>
<tr>
<td>( y_n )</td>
<td>( A_{1n} )</td>
</tr>
</tbody>
</table>
$U_2 \cap SK_a = [(g^{x_1}, g^{y_1}), (g^{x_2}, g^{y_2}), \ldots, (g^{x_n}, g^{y_n}) \cap (g^{a_1}, g^{b_1}), (g^{a_2}, g^{b_2}), \ldots, (g^{a_n}, g^{b_n})].$

If the matching result is not an empty set, then 0 is output. If the matching result is an empty set, then the user attribute private key $SK_a$ matches the allowed access policy $U_1$. The calculation is as follows:

$U_1 \cap SK_a = [(g^{x_1}, g^{y_1}), (g^{x_2}, g^{y_2}), \ldots, (g^{x_n}, g^{y_n}) \cap (g^{a_1}, g^{b_1}), (g^{a_2}, g^{b_2}), \ldots, (g^{a_n}, g^{b_n})].$

When the result is $U_1$, the server outputs 1 and sends the ciphertext $CT$. The server packages the successfully matched attributes to generate matching information $p$.

Dec($CT$, $p$, $B_1$) $\rightarrow m$: The user uses the algorithm to decrypt the ciphertext $CT$ sent by the server. First, the user obtains the matching attribute information from the matching information $p$, calculates the decryption parameter, and generates a random number $r \leftarrow Z_p$.

The calculation is as follow:

$$\hat{e}(g^a, g^b)^l \times k \times \frac{\hat{e}(g^a, B'i)}{\hat{e}(g^{a_1}, B'i, g)} =$$

$$\hat{e}(g^a, g^b)^l \times k \times \frac{\hat{e}(g^a, B'i)}{\hat{e}(g^{a_1}, B'i, g)} =$$

$$\hat{e}(g^a, g^b)^l \times k \times \frac{\hat{e}(g^a, B'i)}{\hat{e}(g^{a_1}, B'i, g)} =$$

$$\hat{e}(g, g)^{abt} \times k \times \frac{\hat{e}(g, B'i)}{\hat{e}(g, g)^{abt}} \hat{e}(B'i, g)^{r} = k.$$

Then, plaintext $m$ is solved by using the symmetric decryption algorithm.

5 Security Analysis

The scheme can ensure data security. The data are encrypted by a symmetric algorithm, and the key $k$ uses public key encryption to obtain ciphertext $C_k$. The keyword trapdoor is randomly encrypted; thus the scheme is statistically IND-KGA secure. In addition, the ciphertext $C_k$ is constructed according to the DBDH assumption, and the proposed scheme is statistically IND-CCA secure under the DBDH assumption.

Theorem 1 Indicates that the proposed scheme is statistically IND-KGA secure.

Proof The keyword trapdoor is unrecognizable.

Setup: Challenger $C$ generates a random number $(a, b) \rightarrow Z_p$, and the public parameters $par = (a, b, g, g^a, g^b, G_1, G_2, \hat{e}, q, H).$

Phase 1: Adversary $A$ selects the keyword set $(w_1, w_2, \ldots, w_n)$ and sends it to Challenger $C$. The Challenger outputs the trapdoor set $(T_{w_1}, T_{w_2}, \ldots, T_{w_n})$ generated by the keyword set and sends it to Adversary $A$.

Challenge: Adversary $A$ selects keywords $w_0$ and $w_1$ that are not in the keyword set in Phase 1 from the keyword space. Challenger $C$ flips a random coin $\mu \in (0, 1)$, runs the algorithm Trap($w, par$), and finally, sends the trapdoor $T_{w_0}$ to Adversary $A$.

Phase 2: Adversary $A$ sends the keyword set to Challenger $C$ again, the same as in Phase 1.

Guess: Adversary $A$ outputs $\mu'$. If $\mu' = \mu$, Adversary $A$ wins the game. The scheme is statistically IND-KGA secure. The keyword trapdoor introduces random numbers during encryption; thus the trapdoors generated by the same keyword are different, which can effectively resist a statistical keyword guessing attack.

Theorem 2 The proposed scheme is statistically IND-CCA secure under the DBDH assumption. Adversary $A$ and Challenger $C$ perform the game as follows.

Proof Adversary $A$ cannot recognize the ciphertext under the DBDH assumption.

Setup: The system is established, the security parameter $\lambda$ is generated, and then the algorithm Setup(1$^\lambda$) is run to obtain the security parameters $par = (a, b, g, g^a, g^b, G_1, G_2, \hat{e}, q, H)$ and the encryption parameter $B$ in the system. Given groups $G_1$, $G_2$, and map $\hat{e} : G_1 \times G_1 \rightarrow G_2$, Challenger $C$ randomly generates $(a, b, c, z) \rightarrow Z_p$ to generate a quintuple $T_0$. $T_0 = (g, A = g^a, B = g^b, C = g^c, Z = \hat{e}(g, g)^z)$.

Phase 1: Adversary $A$ runs encryption algorithm multiple times.

Challenge: Adversary $A$ randomly selects two keywords $m_0$ and $m_1$ from the plaintext space and sends them to Challenger $C$. Challenger $C$ flips a random coin $\mu \in (0, 1)$, runs the algorithm Enc($m, k, w, B, par$) $\rightarrow (CT, \varphi)$, and finally, sends the ciphertext $CT$ to Adversary $A$.

Phase 2: Same as Phase 1.

Guess: Adversary $A$ outputs a plaintext $\mu'$, if $\mu' = \mu'$, Adversary $A$ wins the security game.

According to the above proof, we conclude that our proposed scheme is statistically IND-CCA secure under the DBDH assumption. The ciphertext structure in the scheme is similar to the five-tuple structure in the DBDH assumption and has unrecognizable properties.

6 Performance Analysis

6.1 Theoretical analysis

In theory, three main aspects of the U-ABE scheme
are compared with those of several other schemes: functionality, storage cost, and communication cost. The symbols in the comparison process are defined as follows: $|p|$ indicates the length of the data element in $Z_p$; $|g|$ indicates the length of the data element in $G$; $|g_T|$ indicates the length of the data element in $G_T$; $n_a$ indicates the number of attributes associated with the user; $n_c$ indicates the number of attributes associated with the ciphertext; $n_k$ indicates the number of attributes in the user key; and $n_{a\cdot}$ indicates the number of attributes of the entire system.

6.1.1 Functional analysis

Table 2 compares the functional differences between the U-ABE scheme and the other three schemes. The revocation mechanism of each scheme is immediately revoked. The difference between U-ABE and the other three schemes lies in the revocation direction. The U-ABE scheme has a denial of access policy and can thus be revoked in both directions. The U-ABE scheme adopts the AND mode for access strategy, whereas the other schemes adopt the Tree or the LSSS mode, which consumes less resources and is more efficient.

6.1.2 Storage cost

In Table 3, the U-ABE scheme is compared with schemes in Refs. [13, 18, 20] for storage cost. It is mainly divided into four components for comparison: AA, DO, CS, and U. In the U-ABE scheme, the main role of the AA is to generate attribute tables, encrypt two-way access policy, and assign private keys. The AA mainly stores random numbers and tables. Thus the storage cost of the AA in the U-ABE scheme is calculated as $(2n_a + 1)|p|$. In terms of DO, the main job of the DO in the scheme is to receive the policy ciphertext, generate the index, encrypt the data, and then upload it to the CS. The storage cost of the DO is calculated as $2|p| + |g|$, which is smaller than the storage cost of schemes in Refs. [13, 18, 20]. On the CS side, the main task of the CS in the U-ABE scheme is to receive the ciphertext data uploaded by the DO and the search information and attribute private key uploaded by the U, and then matching calculation is performed between the two. Therefore, the storage cost is calculated as $(n_c + n_k)|g| + 2|g_T|$. Unlike schemes in Refs. [13, 20], the U-ABE scheme reduces the storage cost of the CS. Finally, in the U aspect, the U’s job in the scheme is to receive the attribute private key $SK_a$ from the AA, calculate the attribute private key $SK_a$, generate the trapdoor, and decrypt the ciphertext. Therefore, the U storage cost is calculated as $2n_k + |p|$, which is lower than the schemes in Refs. [13, 18].

6.1.3 Communication cost

Table 4 shows a theoretical analysis of communication cost, which is mainly divided into the data transmission cost of four lines. The first is AA & U. In U-ABE, the AA and U mainly transmit attributes and attribute parameters; therefore, the calculated communication cost is $2n_k + n_k|p|$. Followed by AA and O, in the U-ABE scheme, the plaintext, ciphertext, and encryption parameters of the two-way access policy are mainly transmitted between the AA and the DO, so the

<table>
<thead>
<tr>
<th>Program</th>
<th>Access structure</th>
<th>Revocation method</th>
<th>Safe question</th>
<th>Revoke direction</th>
</tr>
</thead>
<tbody>
<tr>
<td>Scheme in Ref. [18]</td>
<td>LSSS</td>
<td>Immediate</td>
<td>$g$-parallel BDHE</td>
<td>Forward</td>
</tr>
<tr>
<td>ABKS-UR[13]</td>
<td>Tree</td>
<td>Immediate</td>
<td>DBDH</td>
<td>Forward</td>
</tr>
<tr>
<td>AD-KP-ABE[20]</td>
<td>Tree</td>
<td>Immediate</td>
<td>DBDH</td>
<td>Forward</td>
</tr>
<tr>
<td>U-ABE</td>
<td>AND</td>
<td>Immediate</td>
<td>DBDH</td>
<td>Two-way</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Program</th>
<th>AA</th>
<th>DO</th>
<th>CS</th>
<th>U</th>
</tr>
</thead>
<tbody>
<tr>
<td>Scheme in Ref. [18]</td>
<td>$(4 + n_a)</td>
<td>p</td>
<td>$</td>
<td>$(2 + n_a)</td>
</tr>
<tr>
<td>ABKS-UR[13]</td>
<td>$(3 + n_a)</td>
<td>g</td>
<td>+</td>
<td>g_T</td>
</tr>
<tr>
<td>AD-KP-ABE[20]</td>
<td>$2n_a</td>
<td>p</td>
<td>$</td>
<td>$n_c</td>
</tr>
<tr>
<td>U-ABE</td>
<td>$(2n_a + 1)</td>
<td>p</td>
<td>$</td>
<td>$2</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Program</th>
<th>AA&amp;U</th>
<th>AA&amp;DO</th>
<th>CS&amp;U</th>
<th>CS&amp;DO</th>
</tr>
</thead>
<tbody>
<tr>
<td>Scheme in Ref. [18]</td>
<td>$4</td>
<td>g</td>
<td>+ n_k</td>
<td>g</td>
</tr>
<tr>
<td>ABKS-UR[13]</td>
<td>$2</td>
<td>p</td>
<td>+ (n_c + 3)</td>
<td>g_T</td>
</tr>
<tr>
<td>AD-KP-ABE[20]</td>
<td>$n_k</td>
<td>p</td>
<td>$</td>
<td>$n_c</td>
</tr>
<tr>
<td>U-ABE</td>
<td>$2n_k + n_k</td>
<td>p</td>
<td>$</td>
<td>$4n_c +</td>
</tr>
</tbody>
</table>
communication cost can be calculated as $4n_c + |p|$. The comparison of the scheme in this paper with schemes in Refs. [18,20] shows that communication costs are greatly reduced. In the CS and U process, the server in the U-ABE scheme has two data interactions with the data user, mainly the transmission of the trapdoor ciphertext and attribute private key, thus the communication cost is $|p| + (n_k + 1)|g_T| + n_c|g|$, which is superior to the other three schemes. Finally, in the CS and DO process, the DO in the U-ABE scheme unilaterally uploads the ciphertext data to the CS. Thus the communication cost is calculated as $(|g_T| + 1)n_c + (n_c + 1)|p|$.  

6.2 Experiment analysis

The experimental platform is a 64 bits Windows operating system, on an Intel Core(TM) i5-4570 3.20 GHz CPU with 8.00 GB memory. The experimental code is modified and written based on Pairing-Based Cryptography library (PBC), using a super singular curve in Class A, that is $E(F_q) : y^2 = x^3 + x$. Group $G_s$ is subgroup of $E(F_q)$, the order of the group $G_s$ is 160 bits, and the base field is 58 bits. This experiment is conducted with four considerations: encryption time, private key generation overhead, retrieval time, and decryption time. The relationship among the number of attributes, the number of keywords, and the time cost is tested.

6.2.1 Encryption time

Figure 2 shows a comparison of the encryption time of the U-ABE scheme with schemes in Ref. [18], ABKS-UR[13], and AD-KP-ABE[20]. The analysis shows that the U-ABE is superior to the scheme in Ref. [18] and ABKS-UR[13].

6.2.2 Private key generation time

Figure 3 shows a comparison of the private key generation overhead of the U-ABE, the scheme in Ref. [18], and ABKS-UR[13]. Figure 3 indicates that as the number of attributes submitted by users increases, the private key generation time increases linearly. The private key of U-ABE is generated by hashing and exponential operation, which gives this scheme higher computational efficiency than the other two schemes.

6.2.3 Search time

Figure 4 shows the results of a comparison experiment of the retrieval cost for U-ABE, scheme in Ref. [18], and ABKS-UR[13]. The comparison experiment in Fig. 4 is the effect of the number of attributes in the user’s private key on the retrieval time. Also, the keyword is set to 10, because the attribute matching of U-ABE is matched in the form of a set, the time overhead is much smaller than that of the other two schemes. Although U-ABE joins the random number calculation that forms an unrecognizable trapdoor, because only one bilinear calculation and one exponential operation are used in the trapdoor construction process, the calculation amount is
less than that of the other two schemes.

6.2.4 Decryption time

Figure 5 shows a comparison of the decryption time of U-ABE with scheme in Ref. [18] and AD-KP-ABE[20]. As seen from Fig. 5, the decryption time increases as the number of attributes in the user’s private key increases. The analysis shows that U-ABE has obvious advantages over the scheme in Ref. [18] in terms of decryption time. When the number of attributes in the user’s private key reaches 50, the decryption time of U-ABE is less than 1 s.

The decryption time in U-ABE increases as the number of attributes in the private key increases mainly because of the calculation of the decryption parameters, but the decryption algorithm is only a multiplication operation and the decryption time expansion rate is low.

7 Conclusion and Future Work

In this paper, we propose a ciphertext retrieval scheme called U-ABE based on unrecognizable trapdoors. First, we introduce a denial of access policy that can implement two-way revocation to implement the flexible control of access attributes. Moreover, because the access policy is not embedded in the ciphertext, flexible attribute revocation can be completed by modifying the access policy. Second, we introduce the mechanism of the attribute table. Through the reliable AA, an attribute table with random coordinates is created. The attributes of the users and owners are based on Table 1, which reduces the matching calculation time. We finally use the bilinearity of the bilinear pair to construct a one-time trapdoor that realizes the same keyword with different encryption results each time, ensuring keyword privacy.

Through theoretical safety analysis and experimental verification, we prove that our scheme has high security and retrieval efficiency in real data sets. Future work will be performed to improve the solution in terms of AA credibility, accuracy, and efficiency, and to conduct a profound study of trapdoor security issues with the goal of achieving higher security.

References


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